

# On Valuations in Gödel and Nilpotent Minimum Logics

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## The Euler-Klee-Rota lattice-theoretic characteristic

### Valuation

Let  $L$  be a (bounded) distributive lattice whose bottom element is denoted  $\perp$ . A function  $\nu: L \rightarrow \mathbb{R}$  is a **valuation** if it satisfies  $\nu(\perp) = 0$ , and

$$\nu(x) + \nu(y) = \nu(x \vee y) + \nu(x \wedge y)$$

for all  $x, y \in L$ .

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### Lemma

*Every valuation on a finite distributive lattice  $L$  is uniquely determined by its values at the join-irreducibles of  $L$ .*

Recall that  $x \in L$  is *join-irreducible* if it is not the bottom of  $L$ , and  $x = y \vee z$  implies  $x = y$  or  $x = z$  for all  $y, z \in L$ .

## The Euler-Klee-Rota lattice-theoretic characteristic, definition

(V. Klee 1963; G.-C. Rota 1974)

### Euler characteristic

The **Euler characteristic** of a finite distributive lattice  $L$  is the unique valuation  $\chi: L \rightarrow \mathbb{R}$  such that  $\chi(x) = 1$  for any join-irreducible element  $x \in L$ .

## The Euler-Klee-Rota lattice-theoretic characteristic

- Let  $V$  be a set of *vertices*, and let  $P$  be the poset of subsets of  $V$  ordered by inclusion. The collection  $\mathcal{L}$  of lower sets of  $P$  is a (bounded) distributive lattice under  $\cap$ ,  $\cup$ .

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- It turns out that  $\chi$  agrees with the classical Euler characteristic on each simplicial complex  $\Sigma \in \mathcal{L}$ .



## Outline

- 1 Euler Characteristic of a formula in classical propositional logic
- 2 Euler Characteristic of a formula in Gödel logic
- 3 Euler Characteristic of a formula in Nilpotent Minimum logic

## Euler characteristic of a classical formula

- For an integer  $n \geq 0$ , let  $\text{FORM}_n$  denote the set of formulæ in classical (propositional) logic over the atomic propositions  $X_1, \dots, X_n$  and the logical constant  $\perp$  (*falsum*).

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- Then we say that *the Euler characteristic of  $\varphi$  is  $\chi([\varphi]_{\equiv})$ .*

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A **tautology** is a formula  $\alpha$  such that  $\mu(\alpha) = 1$  for every assignment  $\mu$ .

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**Gödel algebras** are Heyting algebras (=Tarski-Lindenbaum algebras of intuitionistic propositional calculus) satisfying the prelinearity axiom

$$(x \rightarrow y) \vee (y \rightarrow x) = \top .$$

They provide the equivalent algebraic semantics of **Gödel logic**. For an integer  $n \geq 0$ , let us write  $\mathcal{G}_n$  for the Tarski-Lindenbaum algebra of Gödel logic over the variables  $X_1, \dots, X_n$ , that is, the algebra  $\text{FORM}_n / \equiv$ , where  $\equiv$  is the logical equivalence between formulæ.



## Euler characteristic of a formula Gödel logic

### Euler characteristic of a formula in Gödel logic

The Euler characteristic of a formula  $\varphi \in \text{FORM}_n$ , written  $\chi(\varphi)$ , is the number  $\chi([\varphi]_{\equiv})$ , where  $\chi$  is the Euler characteristic of the finite distributive lattice  $\mathcal{G}_n$ .

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### Theorem

Fix an integer  $n \geq 1$ . For any formula  $\varphi \in \text{FORM}_n$ , the Euler characteristic  $\chi(\varphi)$  equals the number of **Boolean** assignments  $\mu: \text{FORM}_n \rightarrow [0, 1]$  such that  $\mu(\varphi) = 1$ .

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In the sense given by this result, the characteristic of a formula as defined above is a **classical** notion – it will not distinguish, for instance, classical from non-classical tautologies.

## Gödel $(k + 1)$ -valued logic

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$\mathbb{G}_{k+1}$  is obtained from  $\mathbb{G}_\infty$ , Gödel (infinite-valued) logic recalled above, by restricting assignments to those taking values in the set

$$V_{k+1} = \{0 = \frac{0}{k}, \frac{1}{k}, \dots, \frac{k-1}{k}, \frac{k}{k} = 1\} \subseteq [0, 1],$$

that is, to  $(k + 1)$ -valued assignments.

## Generalised Euler characteristic of a formula in Gödel logic

For a join-irreducible  $g \in \mathcal{G}_n$ , say  $g$  has **height**  $h(g)$  if the (unique) chain of join-irreducibles below  $g$  in  $\mathcal{G}_n$  has cardinality  $h(g)$ .

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### Generalised Euler characteristic

Fix integers  $n, k \geq 1$ . We write  $\chi_k: \mathcal{G}_n \rightarrow \mathbb{R}$  for the unique valuation on  $\mathcal{G}_n$  that satisfies

$$\chi_k(g) = \min \{h(g), k\}$$

for each join-irreducible element  $g \in \mathcal{G}_n$ . Further, if  $\varphi \in \text{FORM}_n$ , we define  $\chi_k(\varphi) = \chi_k([\varphi]_{\equiv})$ .

It turns out that  $\chi_k$  is a “ $k$ -valued characteristic”, as we proceed to show.



## $n$ -equivalence

Our next aim is to relate  $\chi_k$  with (not necessarily Boolean)  $[0, 1]$ -valued assignments. In general, even if  $n = 1$  and the language boils down to  $\{X_1\}$ , there are uncountably many assignments  $\mu: \{X_1\} \rightarrow [0, 1]$ . However, in Gödel logic there is the following important reduction to finiteness.

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### $n$ -equivalence

Fix integers  $n, k \geq 1$ . We say that two  $(k + 1)$ -valued assignments  $\mu$  and  $\nu$  are *equivalent over the first  $n$  variables*, or just  *$n$ -equivalent*, if and only if for all formulæ  $\varphi(X_1, \dots, X_n)$  of  $\mathbb{G}_{k+1}$ ,  $\mu(\varphi) = 1$  if and only if  $\nu(\varphi) = 1$ . The same definition can be given, *mutatis mutandis*, for  $\mathbb{G}_\infty$ .

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where  $P(n, k) = \sum_{i=1}^k \sum_{j=0}^n \binom{n}{j} T(j, i)$ , and

$$T(n, k) = \begin{cases} 1 & \text{if } k = 1, \\ 0 & \text{if } k > n + 1, \\ \sum_{i=1}^n \binom{n}{i} T(n - i, k - 1) & \text{otherwise.} \end{cases}$$

## Table

	k=1	2	3	4	5	6	7
n=1	2	3	3	3	3	3	3
2	4	9	11	11	11	11	11
3	8	27	45	51	51	51	51
4	16	81	191	275	299	299	299
5	32	243	813	1563	2043	2163	2163
6	64	729	3431	8891	14771	18011	18731
7	128	2187	14325	49731	106851	158931	184131
8	256	6561	59231	272675	757019	1407179	1921259
9	512	19683	242973	1468203	5228043	12200883	20214483

The number  $P(n, k)$  of distinct equivalence classes of  $(k + 1)$ -valued assignments over  $n$  variables.

## Main result

### Theorem

Fix integers  $n, k \geq 1$ , and a formula  $\varphi \in \text{FORM}_n$ .

- 1  $\chi_k(\varphi)$  equals the number of  $(k + 1)$ -valued assignments  $\mu: \text{FORM}_n \rightarrow [0, 1]$  such that  $\mu(\varphi) = 1$ , up to  $n$ -equivalence.

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- 2  $\varphi$  is a tautology in  $\mathbb{G}_{k+1}$  if and only if  $\chi_k(\varphi) = P(n, k)$ .
- 3  $\varphi$  is a tautology in  $\mathbb{G}_\infty$  if and only if it is a tautology in  $\mathbb{G}_{n+2}$  if and only if  $\chi_{n+1}(\varphi) = P(n, n + 1)$ .

## Proof of main result

### Lemma 1

Fix integers  $n, k \geq 1$ , let  $x \in \mathcal{G}_n$  and consider the valuation  $\chi_k: \mathcal{G}_n \rightarrow \mathbb{R}$ . Then,  $\chi_k(x)$  equals the number of join-irreducible elements  $g \in \mathcal{G}_n$  such that  $g \leq x$  and  $h(g) \leq k$ .

## Proof of main result

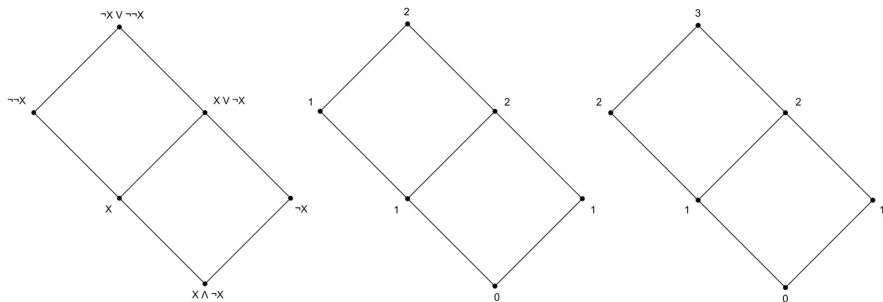
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### Lemma 2

Fix integers  $n, k \geq 1$ , and let  $\varphi \in \text{FORM}_n$ . Let  $O(\varphi, n, k)$  be the set of equivalence classes  $[\mu]_{\equiv_n^k}$  of  $(k+1)$ -valued assignments such that  $\mu(\varphi) = 1$ . Further, let  $J(\varphi, n, k)$  be the set of join-irreducible elements  $x \in \mathcal{G}_n$  such that  $x \leq [\varphi]_{\equiv}$  and  $h(x) \leq k$ . Then there is a bijection between  $O(\varphi, n, k)$  and  $J(\varphi, n, k)$ .

# Example



The Gödel algebra  $\mathcal{G}_1$ , and the values of  $\chi = \chi_1: \mathcal{G}_1 \rightarrow \mathbb{R}$  and  $\chi_2: \mathcal{G}_1 \rightarrow \mathbb{R}$ .

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We can now hope that the Euler characteristic of a formula  $\varphi$  can encode logical information similar to that encoded by the characteristic in the case of Gödel logic.

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We can now hope that the Euler characteristic of a formula  $\varphi$  can encode logical information similar to that encoded by the characteristic in the case of Gödel logic.

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- $[\alpha]_{\equiv}$  is a join irreducible and thus  $\chi(\alpha) = 1$ .

## Positive Euler characteristic of a formula in NM logic

For an element  $x \in \mathcal{NM}_n$ , say  $x$  is **idempotent** if  $x \odot x = x$ .



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### Positive Euler characteristic

We write  $\chi^+ : \mathcal{M}_n \rightarrow \mathbb{R}$  for the unique valuation on  $\mathcal{M}_n$  that satisfies:

- 1**  $\chi^+(x) = 1$  for each idempotent join irreducible element  $x \in \mathcal{M}_n$ .
- 2**  $\chi^+(x \odot x) = \chi^+(x)$  for each  $x \in \mathcal{M}_n$ ;

Further, if  $\varphi \in \text{FORM}_n$ , we define  $\chi^+(\varphi) = \chi^+([\varphi]_{\equiv})$ .

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Note that  $\alpha \odot \alpha = \perp$ , thus  $\chi^+(\alpha) = 0$ .

## Main result

### Theorem

Fix  $n \geq 1$ , and a formula  $\varphi \in \text{FORM}_n$ .

$\chi^+(\varphi)$  equals the number of assignments  $\mu: \text{FORM}_n \rightarrow \{0, \frac{1}{2}, 1\}$  such that  $\mu(\varphi) = 1$ .

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Remark. If  $\varphi$  is a tautology in  $\text{NM}$ , then  $\chi^+(\varphi) = 3^n$ .

## Proof of main result

### Lemma 1

Fix  $n \geq 1$ . Let  $x \in \mathcal{NM}_n$  and consider the valuation  $\chi^+ : \mathcal{NM}_n \rightarrow \mathbb{R}$ . Then,  $\chi^+(x)$  equals the number of minimal idempotent join-irreducible elements  $g \in \mathcal{NM}_n$  such that  $g \leq x$ .

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### Lemma 2

Fix  $n \geq 1$ , and let  $\varphi \in \text{FORM}_n$ . There is a bijection between the set of equivalence classes  $[\mu]_{\equiv_n}$  of assignments to  $\{0, \frac{1}{2}, 1\}$  such that  $\mu(\varphi) = 1$  and the set of minimal idempotent join-irreducible elements  $x \in \mathcal{NM}_n$  such that  $x \leq [\varphi]_{\equiv}$ .

## Further research

- Investigate a **generalised** positive Euler characteristic for NM logic, as done for Gödel logic.
- Investigate the logical content of the Euler characteristic in NM logic.
- Investigate the Euler characteristic in  $NM^{\neg}$ .

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Thank you for your attention.