

Querying with Łukasiewicz logic

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Overview

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Test an intended semantics for Łukasiewicz logic in a real-world situation.

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The Question

What can we ask to our database?

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$$\mu(\alpha \rightarrow \beta) = \min\{1, 1 - (\mu(\alpha) - \mu(\beta))\}$$

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A **tautology** is a formula α such that $\mu(\alpha) = 1$ for every assignment μ .

Łukasiewicz logic in brief (2)

Derived connectives $\top, \vee, \wedge, \leftrightarrow, \oplus, \odot, \ominus$ are defined in the following table, for every formula α and β :

Derived connective	Definition
\top	$\neg \perp$
$\alpha \vee \beta$	$(\alpha \rightarrow \beta) \rightarrow \beta$
$\alpha \wedge \beta$	$\neg(\neg\alpha \vee \neg\beta)$
$\alpha \leftrightarrow \beta$	$(\alpha \rightarrow \beta) \wedge (\beta \rightarrow \alpha)$
$\alpha \oplus \beta$	$\neg\alpha \rightarrow \beta$
$\alpha \odot \beta$	$\neg(\neg\alpha \oplus \neg\beta)$
$\alpha \ominus \beta$	$\neg(\alpha \rightarrow \beta)$

Table : Derived connectives in Łukasiewicz logic.

For each integer $k > 0$ and each formula φ let $1\varphi = \varphi$ and $(k+1)\varphi = \varphi \oplus k\varphi$. Analogously, let $\varphi^1 = \varphi$ and $\varphi^{k+1} = \varphi \odot \varphi^k$.

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- Łukasiewicz logic can deal with properties which has a natural opposite. For instance, **fast** has as its opposite **slow**, with their obvious general meanings, while *red*, the property of being red, might not have a natural opposite.

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- 6 $P \ominus Q$ means "The car is **much more** fast **than** cheap".

The database: data

Our model is a single database table where we have collected cars data. The table contains 4684 records, representing (a subset of) the cars on the Italian market in the year 2014.

Field	Type	Associated Variable
id	int(10) unsigned	-
manufacturer	varchar(50)	-
model	varchar(50)	-
trim	varchar(200)	-
price	int(11)	X0
length	int(11)	X1
width	int(11)	X2
height	int(11)	X3
fuel tank	int(11)	X4
seating capacity	tinyint(4)	X5
car segment	varchar(50)	-
drive	varchar(50)	-
fuel	varchar(50)	-
cubic capacity - cc	int(11)	X6
horsepower	int(11)	X7
power	int(11)	X8
environmental classification	varchar(10)	-
co2 emission	int(11)	X9
gearbox	varchar(50)	-
max speed	smallint(6)	X10
acceleration 0/100	decimal(5,2)	X11
urban cycle consumption	decimal(5,2)	X12
extra-urban cycle consumption	decimal(5,2)	X13
combined cycle consumption	decimal(5,2)	X14

The database: interface

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Łukasiewicz Query:

Connectives:

Symbol	Keyboard	Interpretation
$\neg X$	$\sim, !$	$1 - X$
$X - Y$	$-$	$\max(0, X - Y)$
$X \vee Y$	or	$\max(X, Y)$
$X \wedge Y$	and	$\min(X, Y)$
$X \Rightarrow Y$	\rightarrow	$\min(1, 1 - (X - Y))$
$X \Leftrightarrow Y$	\leftrightarrow	$1 - X - Y $
$X \otimes Y$	$\&ox$	$\max(0, X + Y - 1)$
$X + Y$	$+$	$\min(1, X + Y)$
X^n	X^n	$X \otimes X \otimes \dots \otimes X$
$n \cdot X$	$n * X$	$X + X + \dots + X$
		$n \in \{2, 3, \dots\}$

Record:
 ID: 3003
 Trim: Audi - S1 (8XK)(2014->) - IT - Berl2v3 2.0 TFSI quattro E6 2014 - 2014
 Price: 31500
 Length: 3975
 Width: 1740
 Height: 1417
 Fuel tank: 45

Query:

$Z*((X11^2 \text{ and } (IX12)) - (X0))$

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$2((X11^2 \wedge (\neg X12)) - (X0))$

Results:

Number of true records: 1.

Show Search:

entries

id	allestimento	accelerazione_0_100N	consumo_urbanoN	prezzoN	Results
3003	Audi - S1 (8XK)(2014->) - IT - Berl2v3 2.0 TFSI quattro E6 2014 -	0.82963	0.31875	0.14957	1

Figure : The web page to submit queries

Using linguistic hedges: an example

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The first 10 records of the answer set are:

4272	BMW Serie 3 GT 335d xDrive 2014	[0.793]
4275	BMW Serie 3 GT 335d xDrive 2014	[0.793]
2969	Audi SQ5 (8R) 3.0 TDI DPF quattro 2012	[0.763]
4287	BMW Serie 6 Coupe 640d xDrive 2012	[0.748]
4369	BMW Serie 6 GC 640d xDrive 2013	[0.748]
4433	BMW X4 xDrive 3.5d 2014	[0.748]
4462	BMW Serie 4 Cabrio 435d xDrive 2014	[0.748]
4468	BMW Serie 4 Cabrio 435d xDrive 2014	[0.748]
665	Infiniti Q50 S Hybrid 2013	[0.737]
3028	Audi A7 Sportback 3.0 TDI quattro 2012	[0.733]

Using linguistic hedges: updating queries

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The answer set seems to satisfy our request:

4688	Ford Focus ST EcoBoost 250Cv 2012	[0.556]
3003	Audi S1 2.0 TFSI quattro 2014	[0.551]
2745	Volkswagen Golf VII 2.0 TSI GTI 2013	[0.546]
2205	Renault Megane III Coupe 2.0 TCe 265 2014	[0.544]
4781	Ford Focus Wagon ST 2012	[0.541]
3000	Audi S1 Sportback 2.0 TFSI quattro 2014	[0.538]
478	Seat Leon SC TSI Cupra 2014	[0.537]
3292	Opel Astra J GTC Turbo OPC 2012	[0.535]
496	Seat Leon 2.0 TSI Cupra 2014	[0.530]
2736	Volkswagen Golf VII 2.0 TSI GTI 2013	[0.528]
52	Renault Clio IV 1.6 TCe 200 Monaco GP 2014	[0.526]

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- When we conjunct new conditions, the truth values of the top records of the answer set is decreasing. We shall overcome this obstacle by asking for our final query to be **somewhat satisfied**. For instance, we can rewrite $X_{11}^2 \wedge \neg X_{12} \wedge \neg X_0^3$ as $2(X_{11}^2 \wedge \neg X_{12} \wedge \neg X_0^3)$.

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- It is always possible to rewrite a traditional query in the form $X \leq k$ or $X \geq k$, for $k \in [0, 1]$ a rational value, in a Łukasiewicz query which provide exactly the same result. Nevertheless, this procedure produces very artificial values for exponents and multipliers. For example, $X_{11} \geq 0.875$, and $X_{12} \leq 0.25$ becomes $20(X_{11}^8 \wedge \neg X_{12}^4)$.

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- An assertion like φ can be interpreted as a measure of the level of satisfaction of the buyer: satisfaction begin when pros overcome cons, and is maximal when we have all possible pros, and no cons.

The minus connective \ominus : an example

- The buyer, asked for a definition of a *good car*: "A **good car** has a *remarkable* acceleration, but a *reduced* urban fuel consumption". In symbols: $X_{11}^2 \wedge \neg X_{12}$.

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- The buyer, asked for a definition of a *bad car*: "A **bad car** is an expensive car". In symbols: X_0 .
- In the preceding example, we have *forced* a car to be cheap, using \wedge . We want, instead, to assert that we do not make any difference between a good car with a medium price, and a medium car with a low price. We are asking for a car that is **much more good than bad**.

The minus connective \ominus : an example, cont.

To make such request, we use the following query.

$$X_{11}^2 \wedge \neg X_{12} \ominus \neg X_0^3$$

The minus connective \ominus : an example, cont.

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The answer we obtain querying the database is the following.

3003	Audi S1 2.0 TFSI quattro 2014	[0.510]
3000	Audi S1 Sportback 2.0 TFSI quattro 2014	[0.490]
478	Seat Leon SC 2.0 TSI Cupra 2014	[0.490]
496	Seat Leon 2.0 TSI Cupra 2014	[0.488]
2988	Audi S3 2.0 TFSI quattro 2013	[0.481]
487	Seat Leon SC 2.0 TSI Cupra 2014	[0.480]
4275	BMW Serie 3 GT 335d xDrive 2014	[0.480]
497	Seat Leon 2.0 TSI Cupra 2014	[0.478]
2993	Audi S3 Sportback 2.0 TFSI quattro 2013	[0.476]
1595	BMW Serie 2 Coupe 228i 2014	[0.472]

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- To this end, we need to design the mechanisms of interaction between user and system, and
- ... the system must be equipped with an algorithm able to translate the user's desiderata in a Łukasiewicz formula.

Thank you for your attention.