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Propositional Gödel Logic and Delannoy Paths

Pietro Codara,* Ottavio M. D'Antona, and Vincenzo Marra

Università degli Studi di Milano

*Presenting author

Propositional Gödel Logic

Gödel logic can be semantically defined as the logic of the minimum triangular norm.

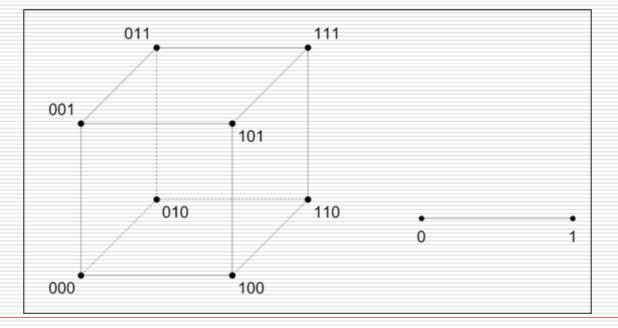
$$x \wedge y = \min(x, y)$$
 $x \rightarrow y = \begin{cases} 1 & \text{if } x \leq y, \\ y & \text{otherwise.} \end{cases}$

$$x \lor y = \max(x, y) \quad \neg x = \begin{cases} 1 & \text{if } x = 0, \\ 0 & \text{otherwise.} \end{cases}$$

Boolean Functions

The algebraic counterpart of the classical propositional logic can be given in terms of Boolean functions, with \land, \lor, \rightarrow , and \neg defined pointwise.

 $f: \{0,1\}^n \to \{0,1\}$



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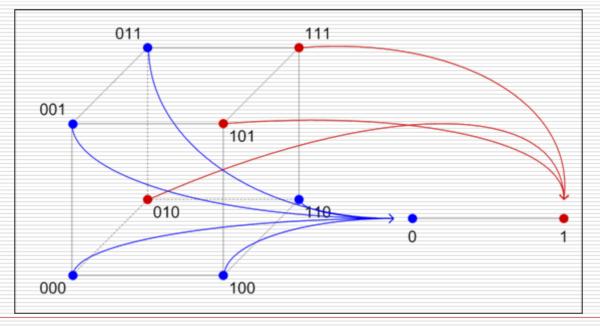
011 111 001 001 001 010 101 0 100

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Enrichment of the *n*-cube

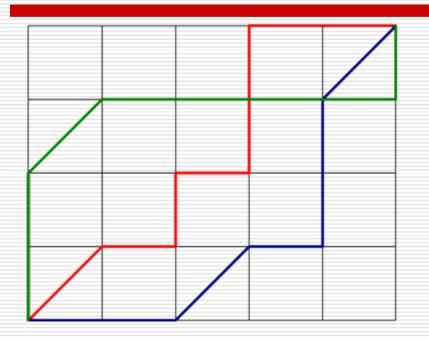
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The answer is given by Delannoy paths.

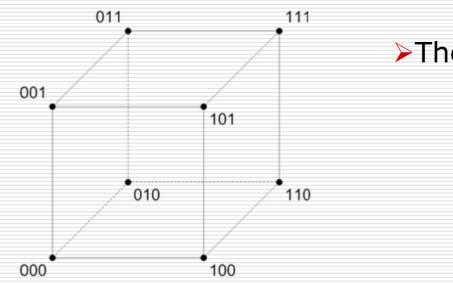
Delannoy Paths





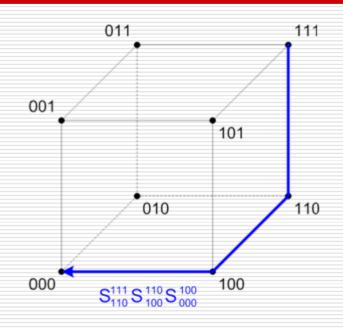
Henri-Auguste Delannoy (1833-1915) French mathematician

□ A Delannoy path is a path in \mathbb{Z}^2 that only uses northward, eastward and northeastward steps.



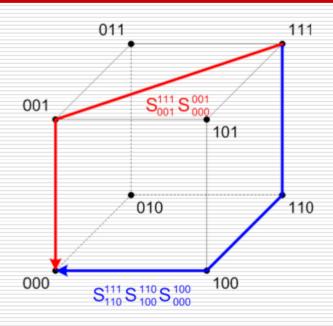
≻The Boolean 3-cube.

- We enrich the Boolean *n*-cube with a variant of Delannoy paths.
- \square We indicate with \mathcal{P}_n the *n*-cube so enriched.



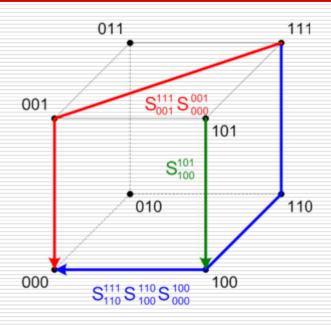
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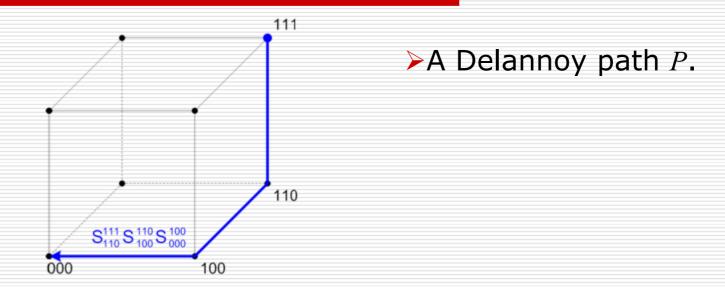
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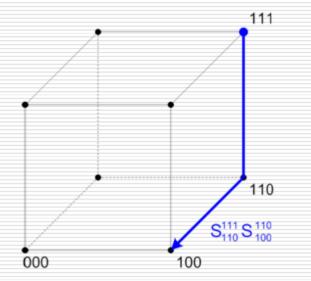


The Boolean 3-cube. A Delannoy path from 111 to 000, in the 3-cube. Another Delannoy path in the 3-cube, with root 111. ➤A Delannoy path with root 101, ending in 100.

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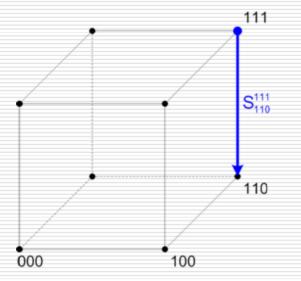


- A subpath of a Delannoy path P is a path having the same root as P, consisting of some initial, consecutive, steps of P (in other words, it is a prefix of P.)
- $\square \quad \Downarrow P$ is the family of all subpaths of *P*.



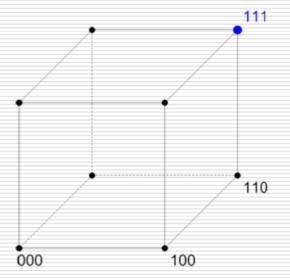
A Delannoy path P. >A subpath of P.

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A Delannoy path *P*.
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➤The shortest subpath of *P*: its root 111.

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From Boolean Functions to Delannoy Functions

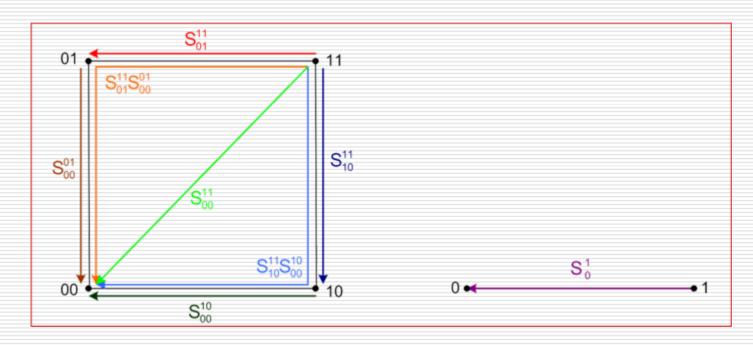
Boolean Functions:

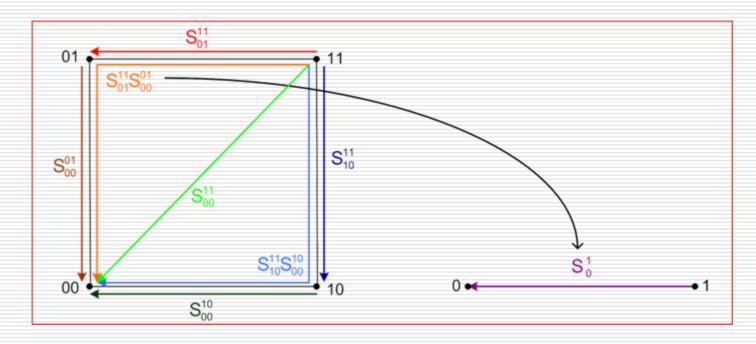
 $f \colon \{0,1\}^n \to \{0,1\}$

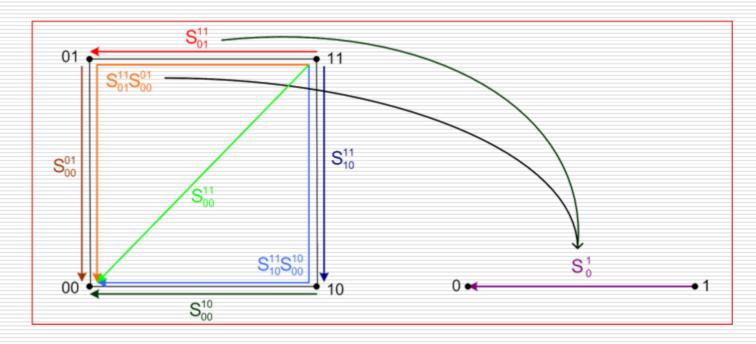
From Boolean Functions to Delannoy Functions

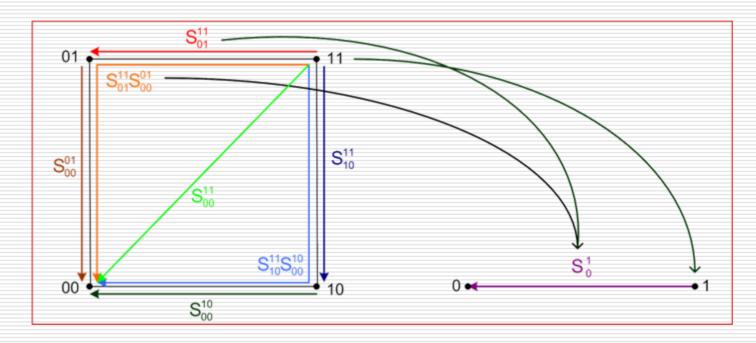
□ Boolean Functions: $f: \{0,1\}^n \rightarrow \{0,1\}$

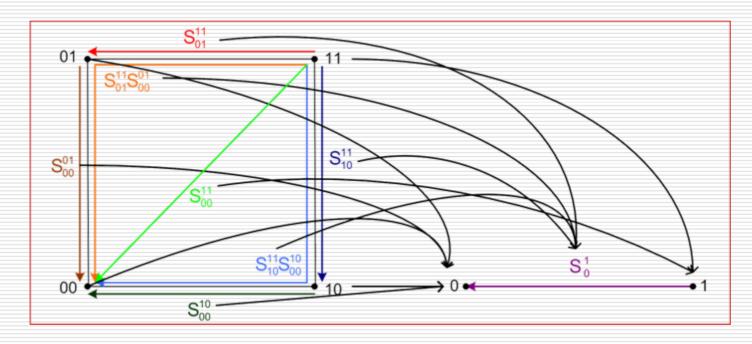
□ Delannoy Functions: $f: \mathcal{P}_n \to \mathcal{P}_1$, with $f(\Downarrow P) = \Downarrow f(P)$.











The Algebra of Delannoy Functions

□ We endow \mathcal{D}_n with appropriate binary operations \land , \lor , \rightarrow , and \neg defined pointwise.

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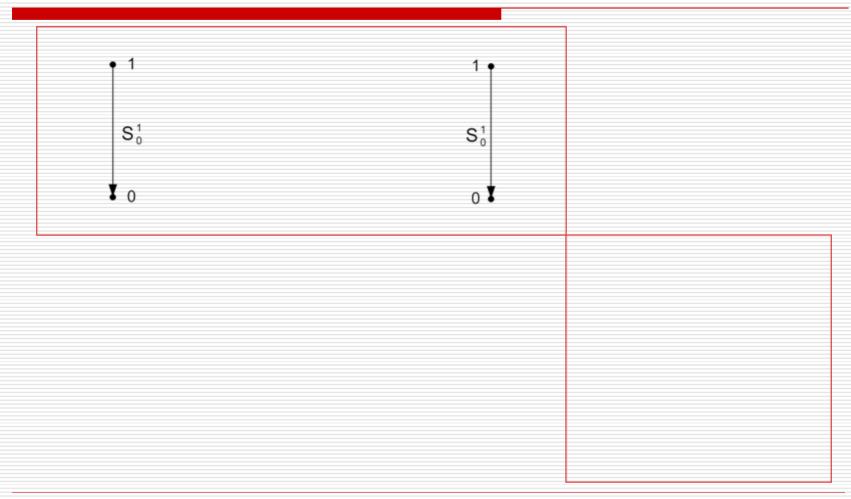
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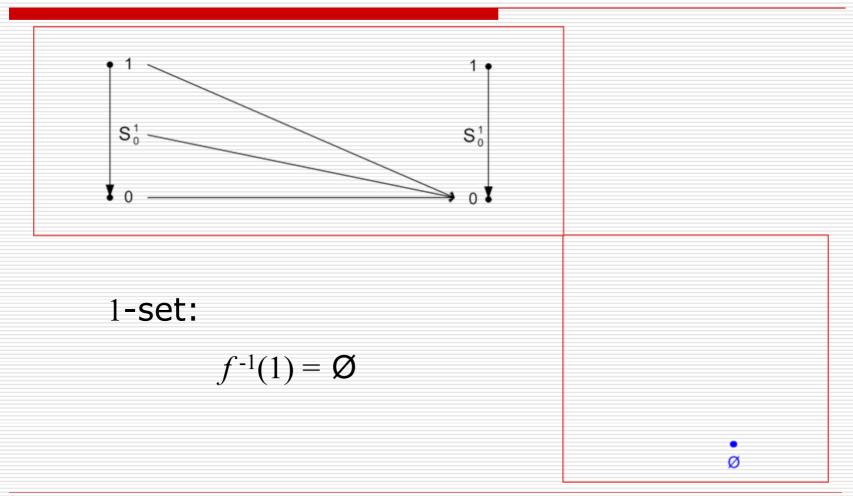
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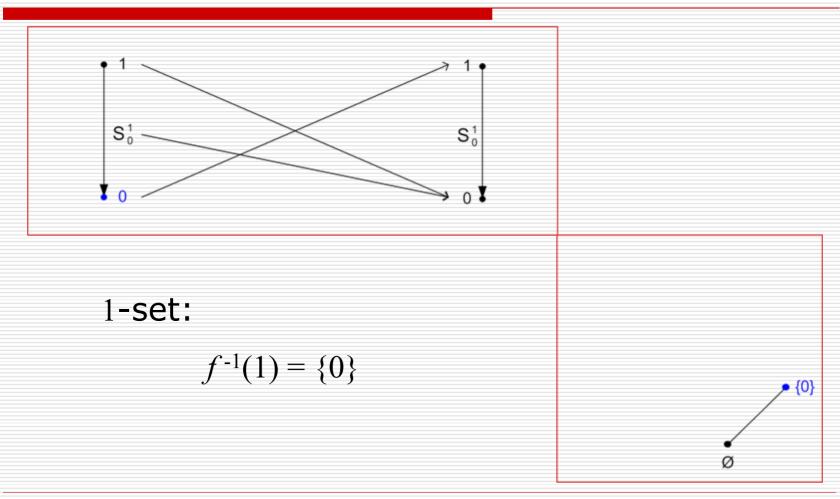
This works because a Delannoy function f is uniquely determined by its 1-set f⁻¹(1), just like in the Boolean case.

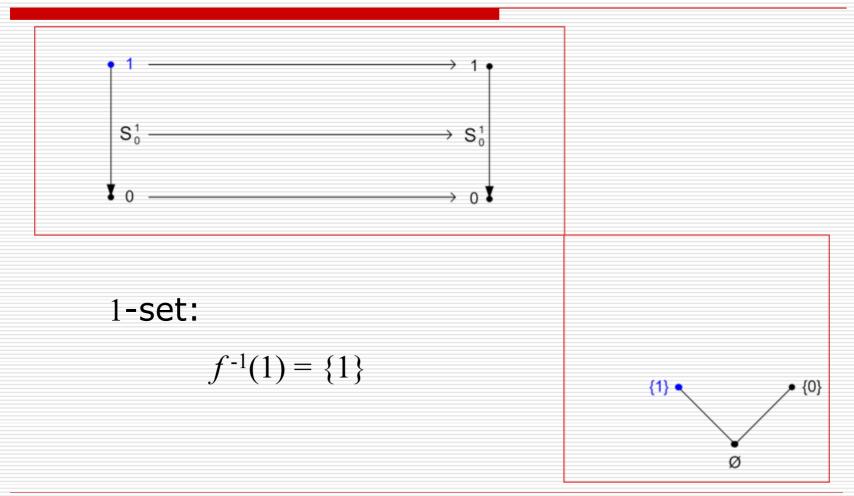
Main Result

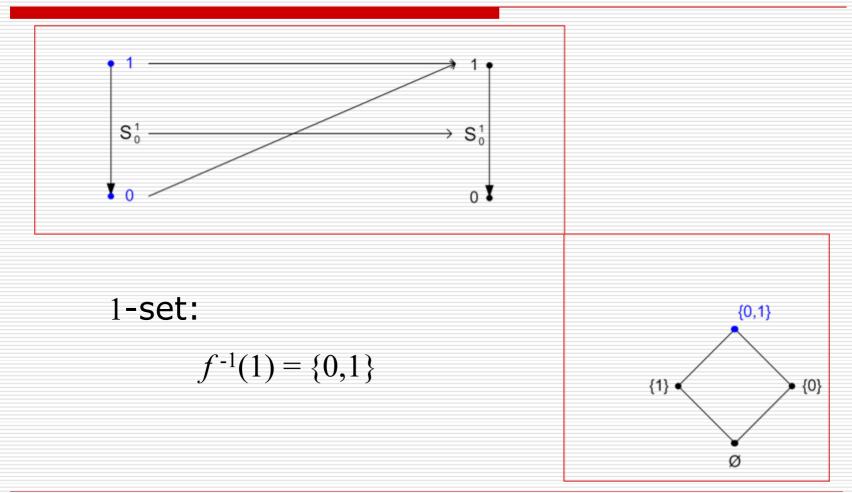
Theorem. $\langle \mathcal{D}_n, \wedge, \vee, \rightarrow, \neg \rangle$ is the free *n*-generated Gödel algebra.

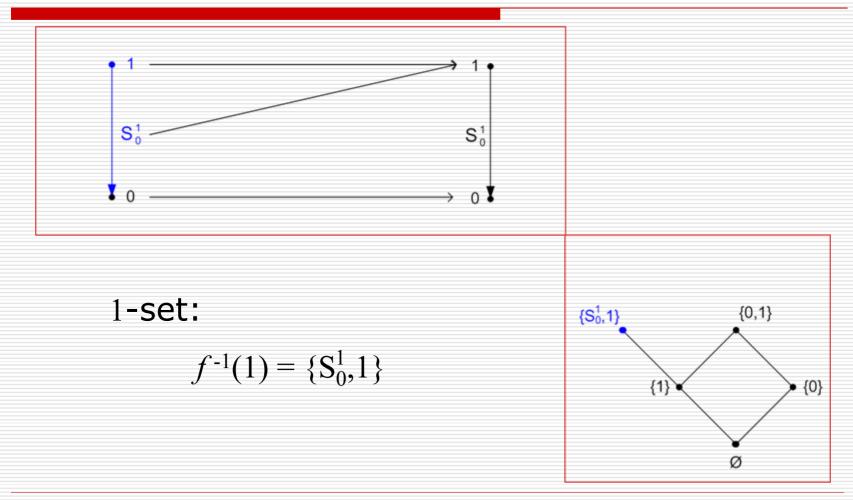


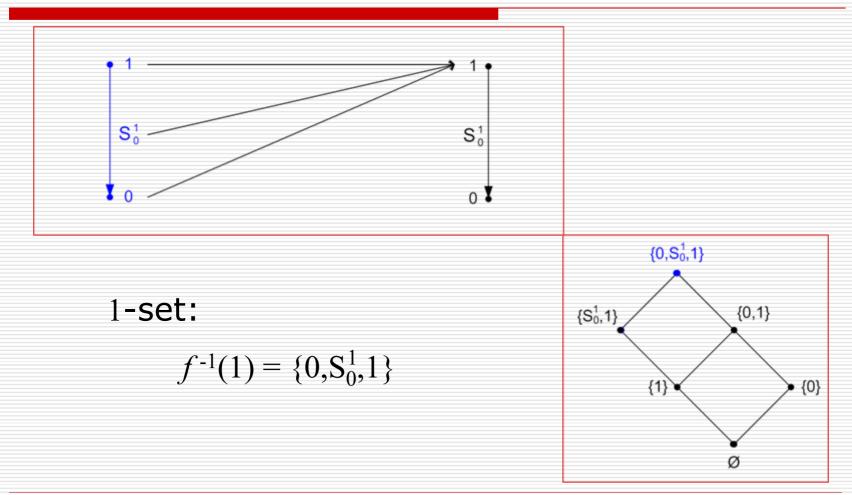












Future Work

D'Antona and Marra, Computing coproducts of finitely presented Gödel algebras, Ann. Pure Appl. Logic, 2006.

Aguzzoli, Gerla, and Marra, Gödel algebras free over finite distributive lattices, preprint. Generalized Delannoy Paths over posets (and not just cubes).

Grazie per l'attenzione

thank you gracias merci gràcies danke obrigado dank u спасибо tack tak teşekkürler