

Combinatorial descriptions of products in the category of forests and open order-preserving maps

Pietro Codara

Dipartimento di Informatica, Università degli Studi di Milano

(joint work with Ottavio M. D'Antona, and Vincenzo Marra)

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Basic notions.

A category of forests

- A **forest** is a finite poset F such that for every $x \in F$, $\downarrow x$ is a chain. A **tree** is a forest with a bottom element.

A category of forests

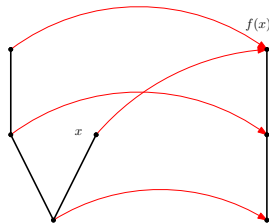
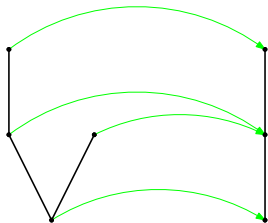
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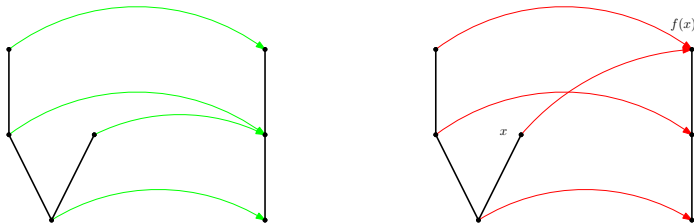
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(The category of forests and open maps is dually equivalent to the category of finitely presented Gödel algebras.)

Ordered partitions, and merged shuffles

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- Let σ and τ be ordered partitions with disjoint supports. An ordered partition θ is a **shuffle** of σ and τ iff σ and τ are subsequences of θ , and $\text{supp}\theta = \text{supp}\sigma \cup \text{supp}\tau$;

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Example. Let $\sigma = \{a|b\}$ and $\tau = \{x\}$. The merged shuffles of σ and τ are: $\{a|b|x\}$, $\{a|x|b\}$, $\{x|a|b\}$, $\{a|bx\}$, $\{ax|b\}$.

Trees of ordered partitions

We can label trees with ordered partitions...

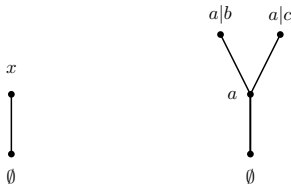
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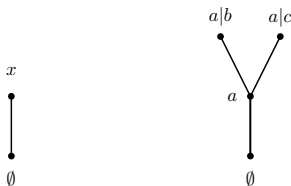
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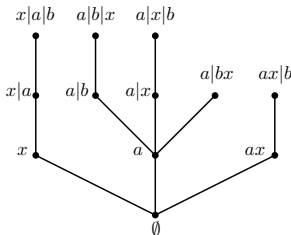
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Product of forests.

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- Let $F = \{T_1, \dots, T_r\}$ and $G = \{U_1, \dots, U_s\}$ be forests.
 $F \times G = \{T_i \times U_j, i \in \{1, \dots, r\}, j \in \{1, \dots, s\}\}.$

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How to compute the product of trees?

Product of trees

Computing the product of trees (an example).

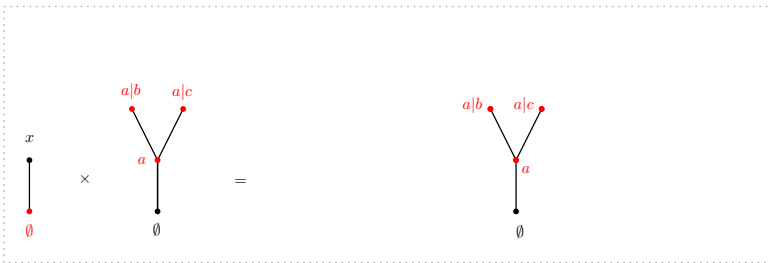
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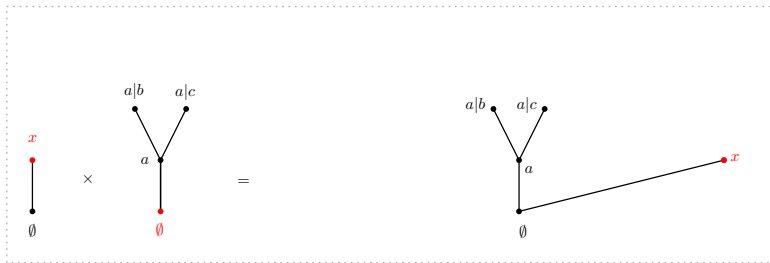
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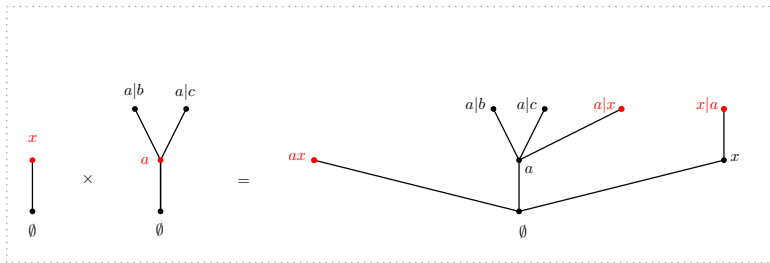
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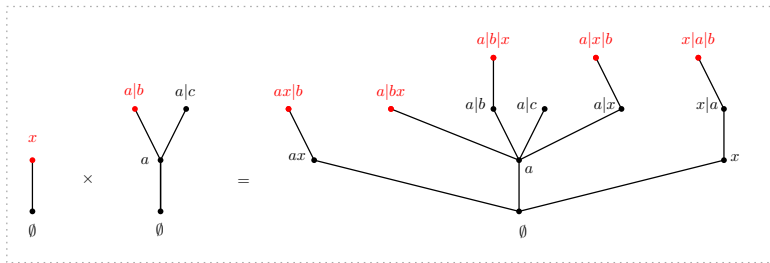
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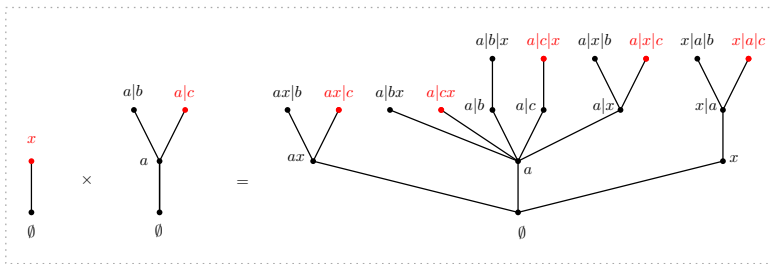
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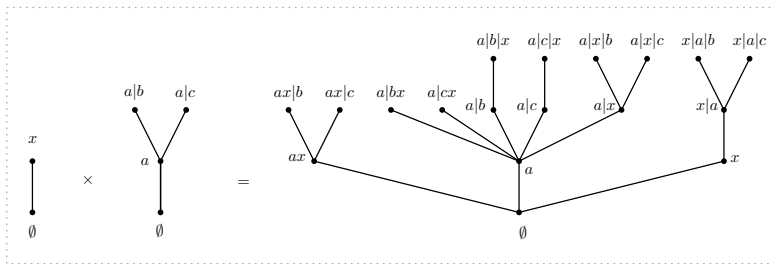
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- $F_{\perp} \times G_{\perp} \cong ((F \times G_{\perp}) + (F \times G) + (F_{\perp} \times G))_{\perp}$.

Product of trees, an alternative way

Computing the product of trees (an example).

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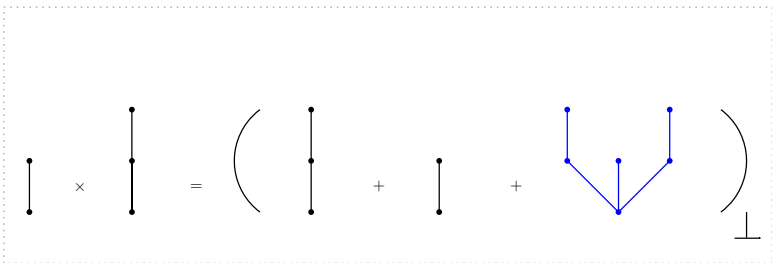
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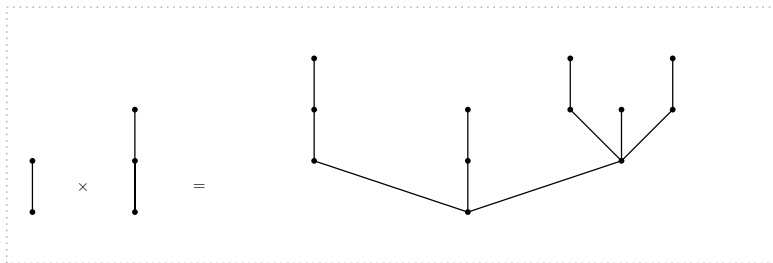
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Enumeration.

Delannoy numbers

The **Delannoy number** $D_{n,m}$ counts the number of lattice paths from $(0,0)$ to (n,m) in which only East, North, and Northeast steps are allowed. Delannoy numbers satisfy the following recurrence relation.

$$D_{n,m} = D_{n-1,m} + D_{n,m-1} + D_{n-1,m-1}$$

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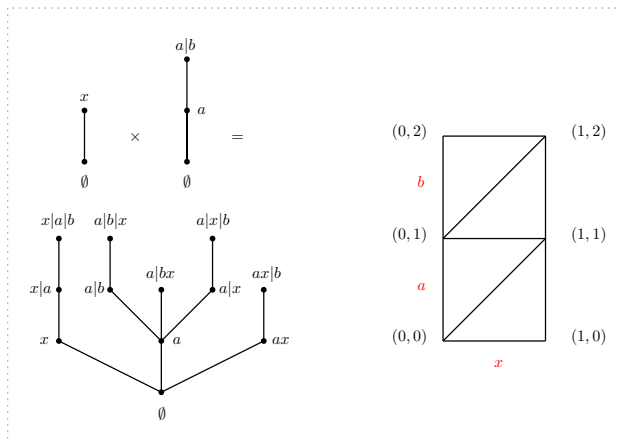
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The following table shows some values of Delannoy numbers.

1	1	1	1	1	1	1	1
1	3	5	7	9	11	13	15
1	5	13	25	41	61	85	113
1	7	25	63	129	231	377	575
1	9	41	129	321	681	1289	2241
1	11	61	231	681	1683	3653	7183

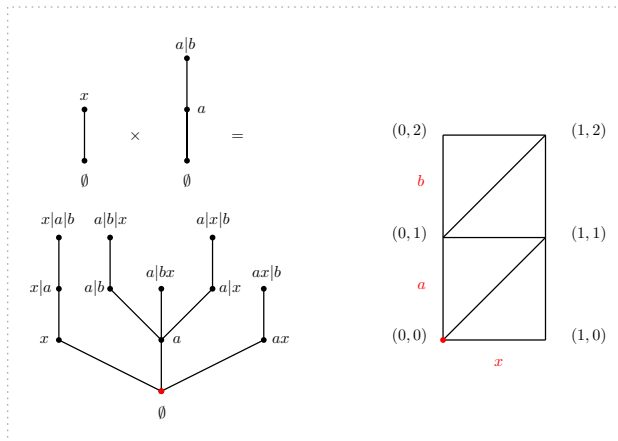
Delannoy numbers and ordered partitions, a simple case

In a simple case, one can associate each element of the product with a Delannoy path, as follows.



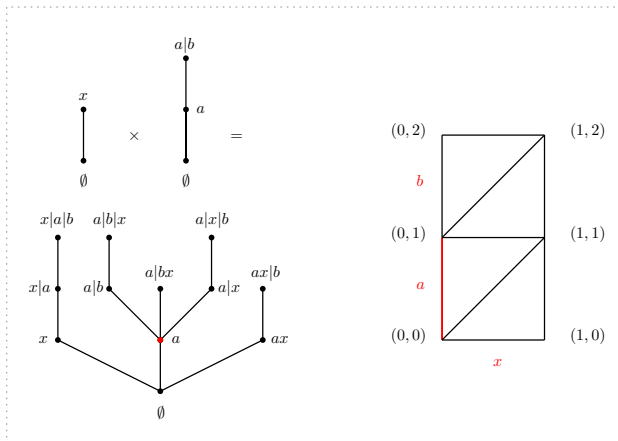
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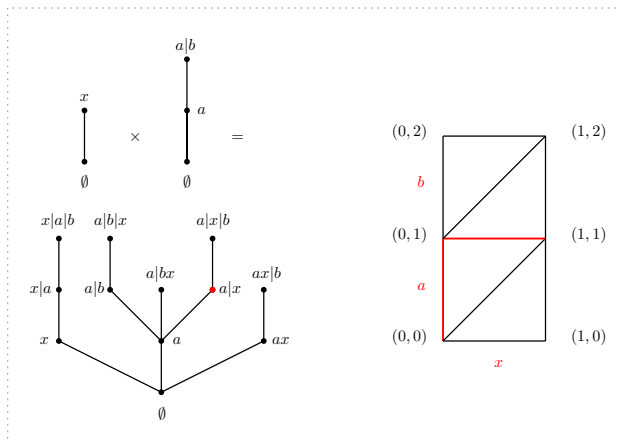
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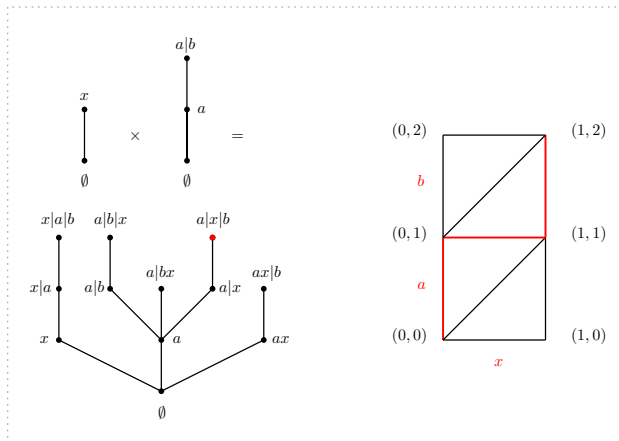
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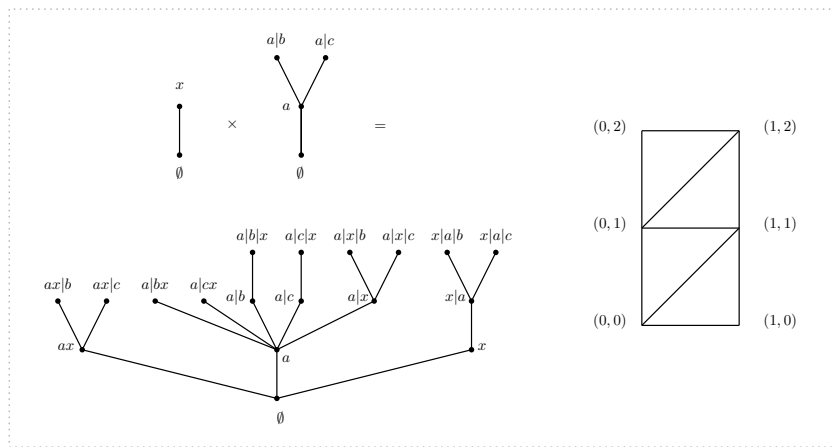
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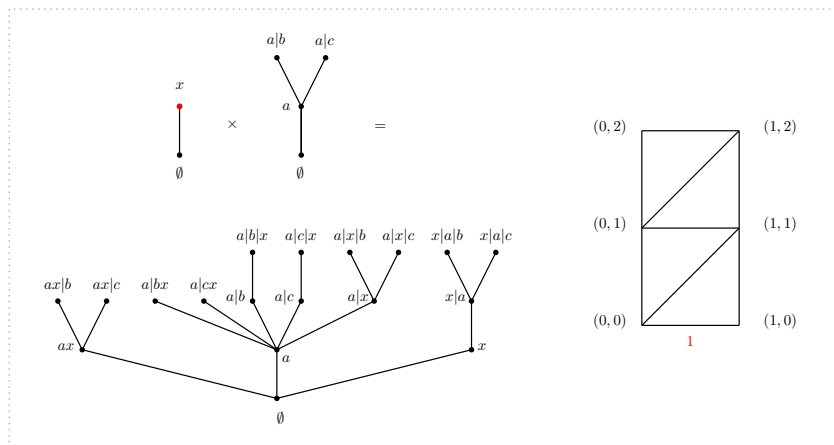
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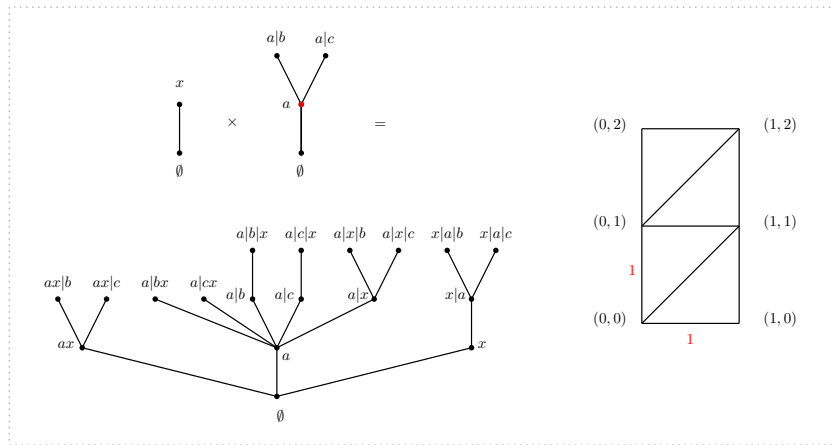
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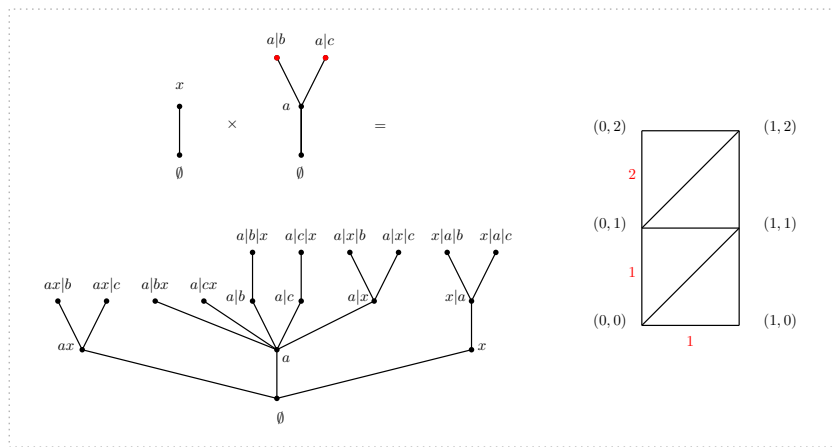
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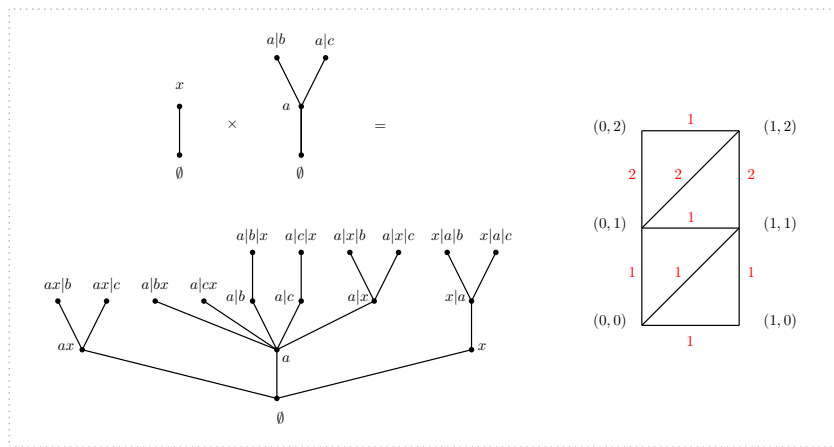
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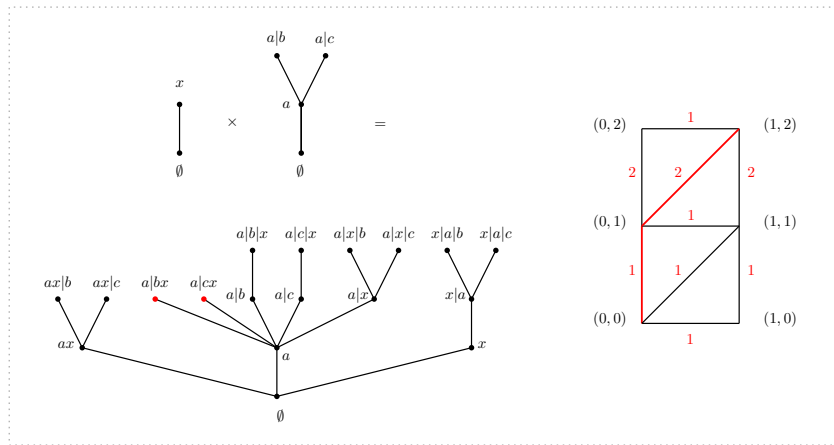
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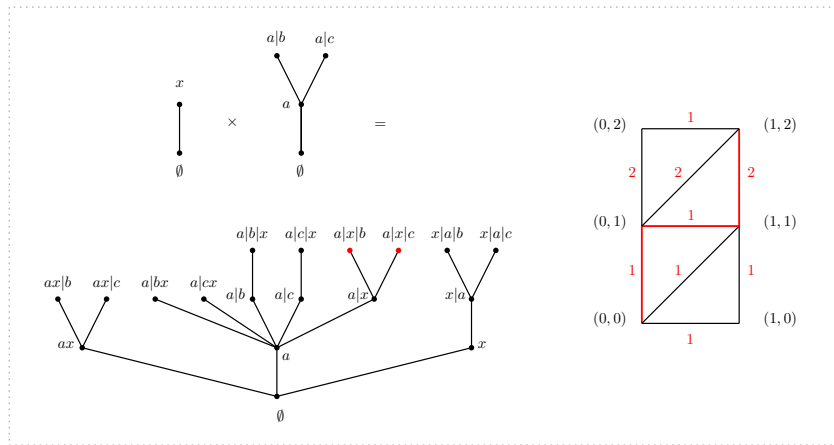
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Formulæ, I

$$|T \times U| = \sum_{i \geq 0} \sum_{j \geq 0} t_i u_j D_{i,j},$$

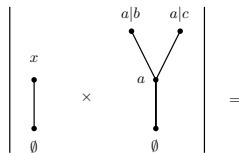
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Example.



$$\begin{aligned}
 &= 1 \cdot 1 \cdot D_{0,0} + 1 \cdot 1 \cdot D_{0,1} + 1 \cdot 2 \cdot D_{0,2} + 1 \cdot 1 \cdot D_{1,0} + 1 \cdot 1 \cdot D_{1,1} + 1 \cdot 2 \cdot D_{1,2} = \\
 &= 1 + 1 + 2 + 1 + 3 + 10 = 18.
 \end{aligned}$$

Formulæ, II

Let L_T be the row vector of length $n = \text{high}(T)$ containing the level numbers of the tree T . Let L_U be the column vector of height $m = \text{high}(U)$ containing the level numbers of the tree U . Let $\mathcal{D}_{n,m}$ be the $n \times m$ matrix such that $\mathcal{D}[i,j] = D_{i,j}$, for each i, j . Then,

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Example.

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 \end{array}$$

Formulae, III

Let $L_{T \times U}[h]$ be h^{th} level number of the product tree $T \times U$,
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Let $L_{T \times U}[h]$ be h^{th} level number of the product tree $T \times U$, for $h \geq 0$. Clearly, $L_{T \times U}[0] = 1$.

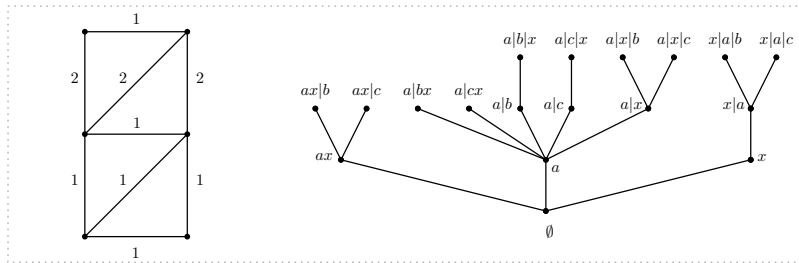
Formulæ, III

Let $L_{T \times U}[h]$ be h^{th} level number of the product tree $T \times U$, for $h \geq 0$. Clearly, $L_{T \times U}[0] = 1$. For $h > 0$, $L_{T \times U}[h]$ is the number of Delannoy paths, with their multiplicity, of length h contained in $(0, 0) - (\text{high}(T), \text{high}(U))$.

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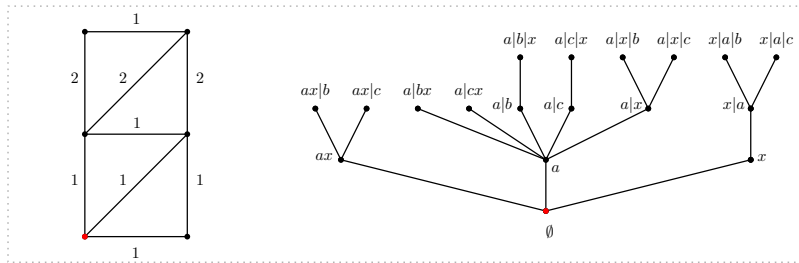
Example.



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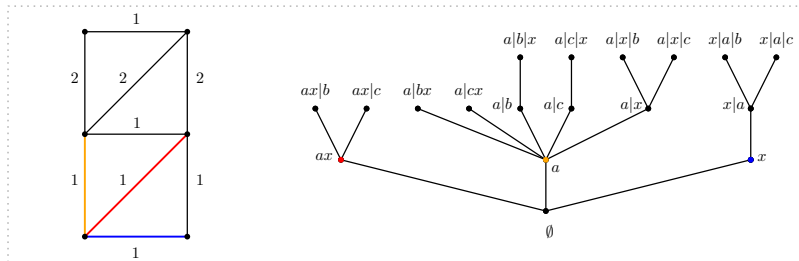
Example.



Formulae, III

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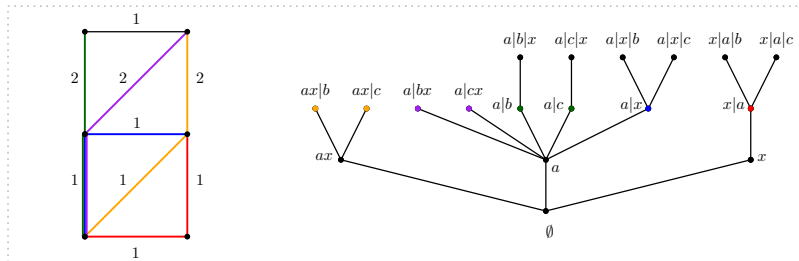
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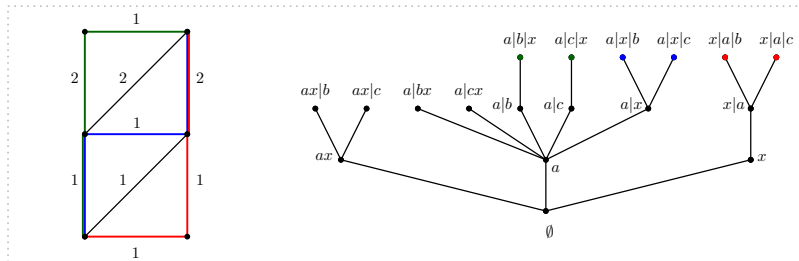
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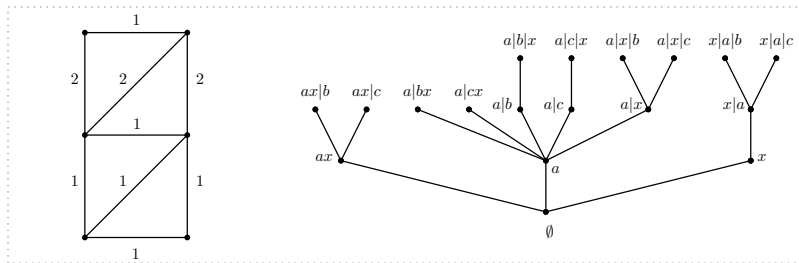
Example.



Formulae, IV

Finally, we can count the number of leaves of the three $T \times U$ after a certain number p of *pruning*, as in the following example.

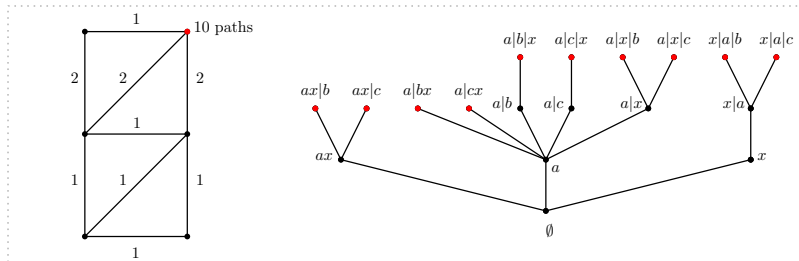
Example.



Formulae, IV

Finally, we can count the number of leaves of the three $T \times U$ after a certain number p of *pruning*, as in the following example.

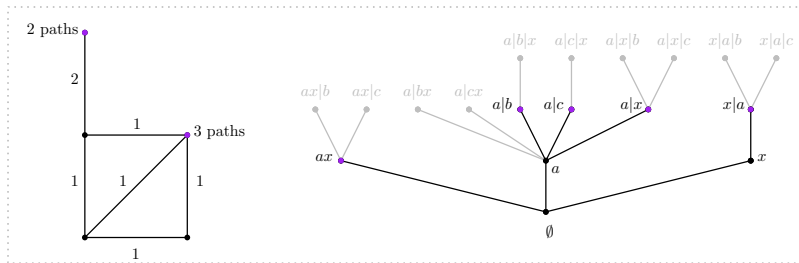
Example.



Formulæ, IV

Finally, we can count the number of leaves of the three $T \times U$ after a certain number p of *pruning*, as in the following example.

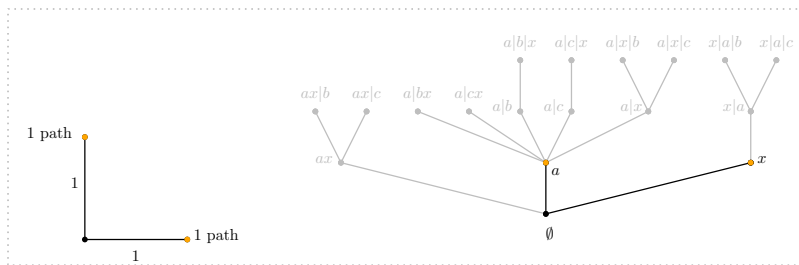
Example.



Formulae, IV

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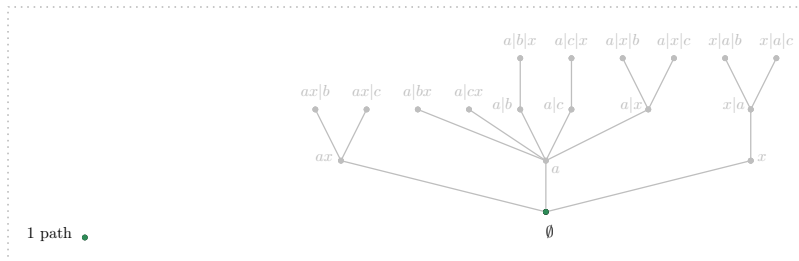
Example.



Formulæ, IV

Finally, we can count the number of leaves of the three $T \times U$ after a certain number p of *pruning*, as in the following example.

Example.



Thank you for your attention.